MB106 QU&NTIT&TIVE TECHNIQUES



MODULE I

LECTURE 12

Transportation Problems-Feasible solution (VAM)

TRANSPORTATION MODEL

Example:

A company has four warehouses and six stores. The warehouses altogether have a surplus of 22 units of a given commodity divided among them as follows:

Warehouses	1	2	3	4
Surplus	5	6	2	9

The six stores altogether need 22 units of the commodity. Individual requirements at stores 1, 2, 3, 4, 5 and 6 are 4, 4, 6, 2, 4 and 2 units respectively.

Cost of shipping one unit of commodity from warehouse I to store j in rupees is given in the matrix below:

	STORES									
		1	2	3	4	5	6			
	1	9	12	9	6	9	10			
WAREHOUSES	2	7	3	7	7	5	5			
	3	6	5	9	11	3	11			
	4	6	8	11	2	2	10			

How should the products be shipped from the warehouses to the stores so that the transportation cost is minimized.

VOGEL'S APPROXIMATION METHOD

Method :

- 1. Note the difference between the lowest and second lowest cost for each row and each column against the row and below the column respectively. This difference shows the penalty incurred due to inability to make an allocation in the least cost cell of that row or column.
- 2. Try to assign maximum possible to the least cost cell of the row or column having maximum penalty.
- Eliminate the row or column whose supply or demand is satisfied completely and recalculate the penalties for the shrunken matrix.
- 4. Repeat steps (i), (ii) and (iii) till all demands are satisfied and supplies exhausted.

	Vogel's Approximation Method:										
	STORES										
			1	2	3	4	5	6			
	WAREHOUSES	1	9	12	9	6	9	10	5 [3]		
		2	7	3	7	7	5	5 2	6/4 [2]		
		3	6	5	9	11	3	11	<mark>2</mark> [2]		
		4	6	8	11	2	2	10	9 [0]		
	Demand		4 [0]	4 [2]	6 [2]	2 [4]	4 [1]	2/0 [5]			
,	Since column 6 has the maximum penalty, maximum allocation is made to the lowest cost cell of this column. Column 6 gets cancelled as demand is fulfilled.										

١	Vogel's Approximation Method:										
	STORES										
			1	2	3	4	5	6			
	WAREHOUSES	1	9	12	9	6	9	10	5 [3]		
		2	7	3	7	7	5	5 2	6/4 [2]		
		3	6	5	9	11	3	11	2 [2]		
		4	6	8	11	2 2	2	10	9 /7 [0]		
	Demand		4 [0]	4 [2]	6 [2]	2/0 [4]	4 [1]	2/0 [5]			
	Since colum lowest cost (n 4 ha cell of	as the ma this colur	aximum p nn. Colun	penalty, n nn 4 gets	naximum cancellec	allocatio I as dema	n is ma nd is fu	ide to the Ifilled.		

Vogel's Approximation Method:									
				STORES				Supply	
		1	2	3	4	5	6		
	1	9	12	9	6	9	10	5 [3][0]	
WAREHUUSES	2	7	3	7	7	5	5 2	6/4 [2]	
	3	6	5	9	11	3	11	2 [2]	
	4	6	8	11	2 2	2 4	10	9 /7/3 [0][4]	
Demand		4 [0]	4 [2]	6 [2]	2/0 [4]	4/0 [1]	2/0 [5]		
Since row 4 has the maximum penalty, maximum allocation is made to the lowest cost cell of this row. Column 5 gets cancelled as demand is fulfilled.									

١	Vogel's Approximation Method:									
					STORES				Supply	
			1	2	3	4	5	6		
	WAREHOUSES	1	9	12	9	6	9	10	5 [3][0]	
		2	7	3 4	7	7	5	5 2	6/4/0 [2][4]	
		3	6	5	9	11	3	11	2 [2][1]	
		4	6	8	11	2 2	2 4	10	9/7/3 [0][4][2]	
	Demand		4 [0]	4/0 [2]	6 [2]	2/0 [4]	4/0 [1]	2/0 [5]		
0	Since row 2 has the maximum penalty, maximum allocation is made to the lowest cost cell of this row. Column 2 and row 2 get cancelled as demand and supply are both fulfilled for them.									

	Jogel's App	roxim	ation M	ethod:					
	1				STORES				Supply
			1	2	3	4	5	6	
		1	9	12	9	6	9	10	5 [3][0]
	WAREHUUSES	2	7	3 4	7	7	5	5 2	6/4/0 [2][4]
		3	6	5	9	11	3	11	2 [2][1][3]
		4	6 3	8	11	2 2	2 4	10	9/7/3/0 [5]
	Demand		4/1 [0]	4/0 [2]	6 [2][0]	2/0 [4]	4/0 [1]	2/0 [5]	
(Since row 4 has the maximum penalty, maximum allocation is made to the lowest cost cell of this row. Row 4 gets cancelled as supply is exhausted.								

V	Vogel's Approximation Method:									
	STORES								Supply	
			1	2	3	4	5	6		
	WAREHOUSES	1	9	12	9	6	9	10	5 [3][0]	
		2	7	3 4	7	7	5	5 2	6/4/0 [2][4]	
		3	6 1	5	9	11	3	11	2/1 [2][1][3]	
		4	6 3	8	11	2 2	2 4	10	9/7/3/0 [5]	
	Demand		4/1/0 [0][3]	4/0 [2]	6 [2][0]	2/0 [4]	4/0 [1]	2/0 [5]		
Ç	Since row 3 and column 1 have the maximum penalty, maximum allocation is made to the lowest cost cell of column 1 and row 3 which intersect at cell 3,1. Column 1 gets cancelled as demand is fulfilled.									

V	'ogel's Approxi	mation	Method:						
			STORES						
			1	2	3	4	5	6	
	WAREHOUSES	1	9	12	<mark>9</mark> 5	6	9	10	5/0 [3][0]
		2	7	3 4	7	7	5	5 2	6/4/0 [2][4]
		3	6 1	5	9 1	11	3	11	2/1/0 [2][1][3]
		4	6 3	8	11	2 2	2 4	10	9/7/3/0 [5]
	Demand		4/1/0 [0][3]	4/0 [2]	6/1 [2][0]	2/0 [4]	4/0 [1]	2/0 [5]	
Sir	Since we are left with only 2 cells, rest of the allocations are made in these cells to satisfy demand and supply restrictions.								
Th	erefore Z=5X9+4>	(3+2X5+2	LX6+1X9+3X6	5+2X2+4X2=1	12				

• TILL WE MEET AGAIN IN THE NEXT CLASS......



